

Home-grown Crypto (Taking a Shiv to a Gun Fight)

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10 AM on Sunday is cruel and unusual punishment. Thank you for coming.

- **Intro**
- Case Studies - Home Grown Crypto Wall of Shame
- Prehash Analysis
- Recommendations and Closing Remarks

- We are not crypto-mathematicians, we're just hackers.
- We regularly see bad crypto.
 - So will you...
- If we can break it, crypto-mathematicians can break it before breakfast.
- We keep seeing these mistakes, so the message bears repeating.

What kinds of things do we see?

- Horrible home-grown “encryption algorithms”
- Giving away the key:
 - For free
 - For a bag of stale Cheetos
 - For a few bucks
- Using good, industry-standard algorithms incorrectly
- Incorrect assumptions about the work factor required to attack the system

- *“Home grown crypto is bad crypto. ... Every programmer tries to build their own encryption algorithm at some point. In one word: Don't.”* Brian Hatch, 2003-01-23
- Security through vigorous handwaving
- Most of the bad crypto we see boils down to security through obscurity, at best

What uses of encryption are we talking about?

- We're not picking on industry-standard algorithms or protocols such as 3DES, AES, or SSL per se ...
- But we will show you cases where we broke implementations protected by them.
- We're talking about using encryption for protecting private data, securing communications, generating session IDs, etc.

- Intro
- Case Studies - Home Grown Crypto Wall of Shame
- Prehash Analysis
- Recommendations and Closing Remarks

- **Case Studies - Home Grown Crypto Wall of Shame**
 - Obfuscation Gone Bad
 - Keys? We Don't Need No Stinking Keys
 - We Don't Need Logic, We've Got Crypto!
 - Opening Your Neighbour's Garage Door
 - We Have Both Kinds: AES and XOR
 - The House Always... Loses?
 - Can't Crack SSL? Just Talk Plaintext!
 - “Take my data. Please!”

Case Study Format

- Description of system or application tested
- Description of flaw or flaws found
 - May include screen shots, graphics, code, etc.
- How we discovered, broke, or reversed it
- How the flaw could have been prevented and/or how we recommended fixing it
 - But only if “fixing it” was a viable option

Obfuscation Gone Bad

- Online Banking Application

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Obfuscation Gone Bad (1)

- Online Banking Application
 - Awarded a seal of approval from a security company that “certifies” security
 - Encrypted account numbers were used to
 - Access account information
 - Download bank statements and cancelled checks
- Unfortunately, the account numbers weren't really encrypted – they were merely obfuscated with a simple substitution cipher.

Obfuscation Gone Bad (2)

- Notice the FNAME value in the following request:
 - TGXPPGGG0001760000250217XRYYYYY
 - It's a sanitized version of the obfuscated account number ...

```
GET /checkimages/GetImage.asp?FNAME=TGXPPGGG0001760000250  
217XRYYYYY&SIDE=FT HTTP/1.0  
Referer: https://somebank.owned.com/checkimages/retrieve  
.asp?TGXPPGGG0001760000250217XRYYYYY  
Connection: Keep-Alive  
Host: somebank.example.com  
[snip]  
Cookie: ASPSESSIONIDGQQGQQZF=PEIAHJKCGBHKNKJIDBEBKNFLG
```

Obfuscation Gone Bad (3)

- The DeEncrypt() routine reveals the ugly truth!
 - Code was obtained from the webserver via a view source bug

```
Sub DeEncrypt(ThisStr,StrLen)
  OldStr = ThisStr
  ThisStr= ""
  For i = 1 to StrLen
    tmpChar = Mid(OldStr,i,1)
    Select Case tmpChar
      Case "G" ThisStr = ThisStr & "0"
      Case "X" ThisStr = ThisStr & "1"
      Case "P" ThisStr = ThisStr & "2"
      Case "K" ThisStr = ThisStr & "3"
      Case "Z" ThisStr = ThisStr & "4"
      Case "L" ThisStr = ThisStr & "5"
      Case "R" ThisStr = ThisStr & "6"
      Case "Y" ThisStr = ThisStr & "7"
      Case "C" ThisStr = ThisStr & "8"
      Case "T" ThisStr = ThisStr & "9"
    End Select
  Next i
[snip]
```

Obfuscation Gone Bad (4)

TGXPPGGG 0001760000250217 XRYYYYYY

90122000 0001760000250217 1677777

- Is this your account number?
- This scheme made it very easy for us to brute force account numbers and access customer statements and cancelled checks

Keys? We Don't Need No Stinking Keys

- Home-grown Single Sign On (SSO) Application

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Keys? We Don't Need No Stinking Keys

- Home-grown Single Sign On (SSO) Application
 - SSO between multiple distributed, un-trusting web servers
 - Usernames and passwords were encrypted and sent as cookies to other sites (in the same DNS domain)
 - Each remote site would decrypt the cookie and pre-populate the login form with the decrypted username
 - The initial login was SSL'ed, but the rest of the applications were not, for “performance reasons”

Keys? We Don't Need No Stinking Keys (2)

- The implementation made knowledge of the crypto and encryption keys unnecessary
- Decrypting any encrypted blob was simply a matter of:
 - Constructing a GET request with the target ciphertext added to a special cookie
 - Submit the request
 - Wait for the server's response
 - View the HTML to recover the plaintext

Keys? We Don't Need No Stinking Keys (3)

- A client logs in; POST arguments include:

```
LOGIN_USERID=foo&LOGIN_PASS=q1w2e3r4
```

- The server sends to the client:

```
Set-Cookie:  
REMEMBERUSERNAME=d47d9e3234c44c3cbc858b43afb1cfb1  
Set-Cookie:  
REMEMBERPASSWORD=cb0d2cf4d49d68cf21634f45a7ec7a1a
```

- If we then send the password blob to the server as a username:

```
GET /loginform.jsp HTTP/1.0  
Cookie: REMEMBERUSERNAME=cb0d2cf4d49d68cf21634f45a7ec7a1a
```

- The server helpfully decrypts it for us:

```
<input class='bodytxt' type="text"  
name="LOGIN_USERID" value='q1w2e3r4'>
```

Keys? We Don't Need No Stinking Keys (4)

- Alone, this flaw was not sufficient to pull off a remote attack (unless you have a nearby victim to sniff). It needed a partner in crime...
- One of the components on one of the SSO'ed webservers was a bulletin-board style app
- Enter XSS – Attack a user via XSS; their encrypted credentials would be exposed to the attacker
- Armed with the ciphertext the attacker could simply issue a modified HTTP request to recover the plaintext

We Don't Need Logic, We've Got Crypto!

- A Web Services application with a Java client

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We Don't Need Logic, We've Got Crypto! (1)

- A Web Services application with a Java client
 - All transactions were SOAP messages – HTTPS POSTs of XML data
 - All client messages used SOAP Digital Signatures
 - The server rejected requests without a valid signature
- We thought we were going to need to do some Java disassembly, extract the private key, inject our own, etc.
 - Turned out it was much, much easier than that

We Don't Need Logic, We've Got Crypto! (2)

- A normal client request might look like:

```
POST /services/FooService HTTP/1.0
Content-Type: text/xml; charset=utf-8
...

<?xml version="1.0" encoding="UTF-8"?>
<soapenv:Envelope ...>
<soapenv:Header>
<wsse:Security ...>
<wsse:UserIsAdmin>no</wsse:UserIsAdmin>
<ds:SignedInfo>
  ...
</ds:SignedInfo>
<ds:SignatureValue>
pRM0+gKSDkpX5v7mWivz78oWDgvGkiOV8vpDNV1VWlVqz2XP68WpydRalI1IYxWzMwtc
aOMkjJHnmRyFBAG92QVe7nUoGIxslMgtEzq+8SKyQu7ldF/xFezMQ0fuf9SjuBCHg==
</ds:SignatureValue>
```

We Don't Need Logic, We've Got Crypto! (3)

- Try to modify the body, and the server will reject it:

```
<?xml version="1.0" encoding="utf-8"?><soapenv:Envelope ...  
<soapenv:Body><soapenv:Fault>  
<faultcode>soapenv:Server.generalException</faultcode>  
<faultstring>WSDoAllReceiver: security processing failed;  
nested exception is:  
    org.apache.ws.security.WSSecurityException: The  
signature verification failed</faultstring>  
<detail> ...
```

- We need to extract the private key the client is using, and sign our own messages. Or do we ?

We Don't Need Logic, We've Got Crypto! (4)

- Although the server checked for valid signatures on client inputs, the client neither verified SSL certificates nor expected the server's responses to be signed
- So... we could forge a message from the server to the client using a standard web-proxy, and then the client would spit it back out to the server at the appropriate time, *complete with a valid digital signature.*

We Don't Need Logic, We've Got Crypto! (5)

- It worked. The server checked that the message had a valid signature – and then blindly trusted the contents!
- There was no actual back-end business logic that checked “Hey wait, is user123 really an admin?”
- This is the Web Services version of a webapp that checks what a user is allowed to do when drawing a menu, but doesn't check when an action is actually performed. Only the oldest business-logic flaw in the book.

Opening Your Neighbour's Garage Door

- Enterprise Security Management Software

Opening Your Neighbour's Garage Door (1)

- Enterprise Security Management Software:
 - Centralized Management Server
 - Management Server pushes configs to and harvests logs from end device nodes
 - Fat GUI client application to administer
 - Thin HTTPS-based client to view logs
 - Client-Server and Server-Node communications are encrypted
 - Fat client and thin client use the same username + password credentials

Opening Your Neighbour's Garage Door (2)

- The Server-Node communications uses AES with some kind of feedback (CBC, etc).
- The encryption keys used are user-configurable in the GUI.
- The GUI only allows alphanumeric keys! It also does not enforce any minimum length, or offer to generate strong keys automatically.
- So, AES or not, it uses a user-supplied encryption key with a reduced keyspace.
- Also, the initial block of an encrypted exchange is always the same, suggesting a fixed IV? Maybe we can dictionary-attack it?
- ...But maybe we don't have to...

Opening Your Neighbour's Garage Door (3)

- The Client-Server communications use SSL.
- The Client software has the server's public key bundled with it, so that it can verify the identity of the server.
- The Server software has a private key bundled with it.
- You read that right. *Every installation of this software, worldwide, has the same private key.*
- All you need is a copy of the server software, and you can extract the private key from it.

Opening Your Neighbour's Garage Door (4)

- As it happens, it's possible to download a demo/trial copy.
- The private key is encrypted in the server installation...
- But of course, since the server needs to know how to decrypt it, the passphrase is conveniently sitting in a .properties file in the same directory as the keystore.
- Once decrypted and extracted, use that RSA private key to ssldump management connections to get administrator's passwords. With that, you can see or do anything; reconfigure the entire security architecture if you want.
- Have you ever deployed an enterprise product that uses SSL, and you didn't have to generate certificates for it? This may apply...

“No Worries, We’re Using 3DES”

- Network-health monitoring appliance

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“No Worries, We’re Using 3DES” (1)

- Network-health monitoring appliance
 - Globally distributed and located in customers' internal networks
 - “Phones home” to appliance vendor to upload data
 - Documentation claimed that sensitive client information was protected through the use of 3DES encryption
- The vendor wanted us to QC their work – kudos to them!
- Unfortunately, the 3DES encrypted keys were being given away

“No Worries, We’re Using 3DES” (2)

- We took a forensic approach:
 - Cracked open the appliance, imaged the hard drive
 - Used FTimes (ftimes.sourceforge.net) to scour the image for clues (credentials, keys, code, DB accounts, etc.)
- A captured transmission revealed that:
 - HTTP version 1.0 and ftp were used to POST data to a vendor site
 - The body of the upload contained binary (3DES?) data that had been encoded as ASCII
- So far so good, right?

“No Worries, We’re Using 3DES” (3)

- We thought so too... until we found the tool that did the encryption:
 - The key was passed in on the command-line
- We traced backwards to find how the key had been generated
- Then, we traced forwards to see how the binary data was generated
- And that's when we saw the key being stuffed in the payload too!

(skip example)

“No Worries, We're Using 3DES” (4)

A capture of the host reporting home looked random. Well, sort of...

```
POST /cgi-bin/app/app_httpdld HTTP/1.0
Host: upload.example.com
Authorization: Basic YWxseW91cmtleXNhcmU6YmVsb25ndG91cwo=
[snip]
Content-Length: 2441
Content-Type: multipart/form-data
```

105860100420976322022480.StatusDownload.rcv

_dftgeqaSay

;<CE[DCDCBFh@h3h>h_^cfd^_^b`^geda`^^`bf^h#D1DEECr?G><?14hDB1>C@?

BDE[851<D8E[C165\D1Bh?EED213;C5BF5B

```
8[y@`z6<^g^T`&"XfnEl_U]4>RZ79tL2ei[3Z*znZBz}}FhnAkq$j[_I_\gvzqJ3X)4U0;)
j`axBe8yLCw=xrx"4apa1.;H1VqA]w3lV;"")+7YeJ9c"e_h!]<1:,5a[+FaA;wQC^w{:s
"I;,(:6pQ\LU7+L}2X\R>21i4c8XQh<]0T$_EgzhLYp7Li@R9&z&5R2"3?x8HY2i&#S6R[
A1@=XdsvbBIZRn0/#\X'3wjVn{R5`ABR\Hwg{<W@[hCR<si*b2BQE4@dV[VVNg:*XE2B:4
:$32Y/VNFar)`|G7|/noDr9}5nD]4}
```

[snip]

“No Worries, We’re Using 3DES” (5)

The encoded stream is also sent via ftp to an upload server.

```
220 FTP server ready.
USER upload2
331 Password required for upload2..
PASS L33tDuD3
230 User upload2 logged in. Access restrictions apply.
[snip]
STOR 238051003320976321714736.DailyDownload.rcv.
150 Opening BINARY mode data connection for
    238051003320976321714736.DailyDownload.rcv.
`sfrbdgqS_d^y
;<CE[DCDCBFh@hJh4`h`af^c_^^aa`^geda`_e_beadhr19<Ir?G><?14hDB1>C@?BDE[C165\D1Bh?
EED213;C5BF5B
89UvrdN=Y!I<d9r4a<E]SHinL;_!$!pJg[*B(sDU|p,X`vF0}/e_s9fU@b[K{DdW=Iqj1B[/q CU2IR
xQF0Tz|oror2mL[E&if'V)k!4I(|hVyT3NJsh05s"nj2E/i$;o`[`LcSXV1E6/qQy1'S!0V\SvpYI1
(m,SF{Hnq.rqh#)l!k@,$6k> 1;;Yk+T'h3a4nWp6(2\#J35h_C/*x9 )?q[mUwed"$A9?Z1I^#Q,!
IYqZn`3 jW4.|Xe:5G67:z@'!.!x|3s0UhK#v!\ ,#=#|bINW>d+6"IL2@(c#WQG14Y8h1,z_424`}F.
[snip]
```

“No Worries, We’re Using 3DES” (6)

```
#!/usr/bin/perl -wT
use strict;
$::map = <<'EOMAP';
N {" {# {$ {% {& {' {( {) ... [snip]
EOMAP
@::toenc = split(/ /, $::map);
my @enc_regex;
for (my $i=0; $i<$#:toenc; $i++) {
    $::toasc{ $::toenc[$i] } = $i;
    my $token = $::toenc[$i];
    $token =~ s/([!]?{ }.+*()&|\|^\\|)\|^$1/g;
    push(@enc_regex, $token);
};
$::regex = join('|', @enc_regex);
while(@ARGV and $ARGV[0] =~ /^-([a-zA-Z])/) {
    my $arg = $1;
    die "Usage: $0 [-h] [-x] [-r] [-l string]\n" if ($arg eq 'h');
    $::func = \&Decode if ($arg eq 'r');
    $::hexit = 1 if ($arg eq 'X');
    shift;
};
$::func = \&Encode unless ($::func);
while(<>) {
    &$::func($_);
};
sub Decode {
    while (m/($::regex)/g) {
        print chr($::toasc{$1});
    };
};
sub Encode {
    foreach my $char (split(//,$_)) {
        my $ascii = ord($char);
        my $mapped = $::toenc[$ascii];
[snip]
```

The ASCII encoding wasn't anything standard (uuencode, base64, hex-encoding, etc.). But we had their binary.

After we made a quick dictionary of all possible byte values, a quick perl equivalent encoder / decoder was produced.

“No Worries, We’re Using 3DES” (7)

Decoding the headers provided us with clues.

The three header lines contain transfer metadata: filename, appliance name, transfer type, various control flags, and oh... the 3DES encryption keys.

```
359574702420976322006928.DailyDownload.rcv
tbae_ofeS_`dy
;<CE[DCDCBFh@h3h>hacgcebe^`b`^geda``^^dg`fhr19<Ir?G><?
14hDB1>C@?BDE[C165\D1Bh?EED213;C5BF5B
```

Decoded:

```
359574702420976322006928.DailyDownload.rcv
F4371A87 126K
kls-tstsrv:p:c:n:359574702420976322006928:DailyDownload:
transport-vault.tar:outbackserver
```

“No Worries, We're Using 3DES” (8)

Example decoding: putting all the pieces together

Header lines extracted from a transfer:

```
238051003320976321714736.DailyDownload.rcv
`sfrbdgqS_d^y
;<CE[DCDCBFh@hJh4`h`af^c_^^aa`^geda`_e_beadhr19<Ir?G><?
 14hDB1>C@?BDE[C165\D1Bh?EED213;C5BF5B
```

Decoded Header lines:

```
238051003320976321714736.DailyDownload.rcv
2E8D469C 160K
kls-tstsrv:p:z:d2:238051003320976321714736:DailyDownload:
  transport-vault.tar:outbackserver
```

Decode the body using our script:

```
$ encode_char.pl -r <body.tar.enc.hdr.gz.3des.enc
  >body.tar.enc.hdr.gz.3des
```


“No Worries, We’re Using 3DES” (9)

Example decoding: putting all the pieces together, cont

The key is comprised of the transaction ID + HOSTNAME:

Convert the key to hexadecimal and crop to 32 characters:

The hexadecimal value of the ID always exceeds 32 characters, making the inclusion of the HOSTNAME in the key irrelevant.

```
echo '238051003320976321714736kls-tstsr' | encode_char.pl -X | cut -c 1-32  
6061665E635F5E5E6161605E67656461
```

Decrypted the data using the key, and ‘des’ from the ancient libdes package:

```
des -d -3 -h -k 6061665E635F5E5E6161605E67656461 body.tar.enc.hdr.gz.3des  
body.tar.enc.hdr.gz
```

Decompress the file:

```
gzip -d body.tar.enc.hdr.gz
```

Remove a redundant header at the end of the file:

```
perl -ne 'print $last; $last=$_' <body.tar.enc.hdr >body.tar.enc
```

Decode the charmapped encoded file:

```
encode_char.pl -r <body.tar.enc >body.tar
```

“No Worries, We’re Using 3DES” (10)

- Here's a sample payload after it had been cracked

```
tar -tvf body.tar
-rw-rw-r-- tuser/users      0 2004-05-14 15:08 out/logs.download
-rw-rw-r-- tuser/users      0 2004-05-14 15:08 out/tuserreportrundetailsvault.download
-rw-rw-r-- tuser/users      0 2004-05-14 15:08 out/tuserbxmcardresourcevault.download
-rw-rw-r-- tuser/users      0 2004-05-14 15:08 out/tuserbxmchannelusagevault.download
-rw-rw-r-- tuser/users    571 2004-05-14 15:08 out/modulesvault.download
-rw-rw-r-- tuser/users    924 2004-05-14 15:08 out/modulesattrvault.download
-rw-rw-r-- tuser/users    143 2004-05-14 15:08 out/downloadmastervault.download
-rw-rw-r-- tuser/users   5357 2004-05-14 15:08 out/schedulervault.download
-rw-rw-r-- tuser/users      0 2004-05-14 15:08 out/dailyrvault.download
-rw-r----- tuser/users 187204 2004-05-14 15:08 out/logs/seedfmgr.log
-rw-r----- tuser/users 1047963 2004-05-14 15:08 out/logs/scheduler.log
[snip]
```

“No Worries, We’re Using 3DES” (11)

- The information was encrypted with 3DES prior to transfer (as claimed in the documentation)
- It was then ASCII encoded with a substitution cipher
- BUT... the first few bytes of the transfer contained the encryption key used to encrypt the rest of the file!
- The vendor didn’t mention that... :/

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We Have Both Kinds: AES and XOR

- “Convergent technologies” application

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We Have Both Kinds: AES and XOR (1)

- “Convergent technologies” application:
 - Fat MS Windows client
 - Clients for both x86 W2K, and WinCE
 - Mixture of web-based and proprietary communications
 - AES encryption used in some places and XOR encoding used in others
- Both had problems...

We Have Both Kinds: AES and XOR (2)

- Initially, we found the secret AES key in a DLL using traditional reversing techniques
 - Strings, hex dumps, static disassembly of subtle routines like: `Encrypt()` and `Decrypt()`
 - Umm, great key guys: “1234567890abcdef”
- Later, we used FTimes in dig mode to search all files for the observed key pattern:
 - `DigStringRegExp=\x00[0-9A-Fa-f]{16}\x00`
 - This often reveals additional keys, but not this time

We Have Both Kinds: AES and XOR (3)

- Their choice of key suggested that the total key space was severely limited:
 - Total key space = 256^{16} or 2^{128}
 - 340,282,366,920,938,463,463,374,607,431,768,211,456
 - Hex characters only key space = 16^{16} or 2^{64}
 - 18,446,744,073,709,551,616
- The difference is huge!
- But is the crippled key space big enough?
 - Are there other limiters in play?

We Have Both Kinds: AES and XOR (4)

- Initially, our plan was to use the recovered key and an alternate AES implementation to handle our encryption/decryption needs
- But the exported encrypt/decrypt routines were too enticing, so we chose to abuse them instead
 - Eliminated potential implementation gotchas
 - Required disassembly to determine call setup
 - But the results were worth the effort...

We Have Both Kinds: AES and XOR (5)

- Here's a pseudo C implementation that shows how we used the exported Decrypt routine

```
...
#define KEY "1234567890abcdef"
typedef DWORD(__stdcall *MYROUTINE)(LPTSTR, LPTSTR, DWORD, LPTSTR);
int main() {
    ...
    Handle = LoadLibrary("Application.dll");
    if (Handle != NULL) {
        Routine = (MYROUTINE) GetProcAddress(Handle, "Decrypt");
        if (Routine != NULL) {
            Routine(CipherText, PlainText, CHUNK_SIZE, Key);
        }
    }
    ...
}
```

We Have Both Kinds: AES and XOR (6)

- Now we could encrypt and/or decrypt
 - Local files (e.g., user profile)
 - Messages transmitted to/from the server
- This allowed us to craft custom files and messages
- It also revealed a new problem:
 - The implementation used ECB instead of CBC
 - ECB – Electronic Code Book
 - CBC – Cipher Block Chaining
 - ECB was not a good choice in this case... (skip example)

We Have Both Kinds: AES and XOR (7)

```
00000000 c5 e1 c6 6f 93 4a 73 f1 5a 49 e5 27 fe d9 b6 58 |...o.Js.ZI.'...X|
00000010 e3 6f 87 24 f3 b1 ed XX XX XX XX XX XX XX XX XX |.o.$.....|
00000020 ef 0b f2 b7 c9 0f b5 e9 80 de 51 9a 16 f8 82 40 |.....Q....@|
00000030 f0 60 a8 d5 35 74 84 1c e5 97 43 XX XX XX XX XX |.`..5t....C....|
00000040 ff 5b 0c e9 07 8b d0 8e d2 fe a4 0e fb a5 ca 5d |.[.....]|
```

```
00000000 3c 43 75 73 74 6f 6d 65 72 3e 3c 41 63 63 6f 75 |<Customer><Accou|
00000010 6e 74 4e 61 6d 65 3e XX XX XX XX XX XX XX XX XX |ntName>XXXXXXXXXX|
00000020 34 3c 2f 41 63 63 6f 75 6e 74 4e 61 6d 65 3e 3c |4</AccountName><|
00000030 47 72 6f 75 70 73 3e 3c 49 44 3e XX XX XX XX XX |Groups><ID>XXXXXX|
00000040 3c 2f 49 44 3e 3c 4e 61 6d 65 3e 46 72 69 65 6e |</ID><Name>Frien|
```

We Have Both Kinds: AES and XOR (8)

```
00000000 c4 e1 c6 6f 93 4a 73 f1 5a 49 e5 27 fe d9 b6 58 |...o.Js.ZI.'...X|
00000010 e3 6f 87 24 f3 b1 ed XX XX XX XX XX XX XX XX |.o.$.....|
00000020 ef 0b f2 b7 c9 0f b5 e9 80 de 51 9a 16 f8 82 40 |.....Q....@|
00000030 f0 60 a8 d5 35 74 84 1c e5 97 43 XX XX XX XX XX |.`..5t....C....|
00000040 ff 5b 0c e9 07 8b d0 8e d2 fe a4 0e fb a5 ca 5d |.[.....]|
```

```
00000000 71 f7 95 1e 19 d3 bd 5e 6d 71 e9 eb 45 18 61 84 |q.....^mq..E.a.|
00000010 6e 74 4e 61 6d 65 3e XX XX XX XX XX XX XX XX |ntName>XXXXXXXXXX|
00000020 34 3c 2f 41 63 63 6f 75 6e 74 4e 61 6d 65 3e 3c |4</AccountName><|
00000030 47 72 6f 75 70 73 3e 3c 49 44 3e XX XX XX XX XX |Groups><ID>XXXXXX|
00000040 3c 2f 49 44 3e 3c 4e 61 6d 65 3e 46 72 69 65 6e |</ID><Name>Frien|
```

We Have Both Kinds: AES and XOR (9)

```
00000000 11 11 11 11 11 11 11 11 11 11 11 11 11 11 11 | ..... |
00000010 e3 6f 87 24 f3 b1 ed XX XX XX XX XX XX XX XX | .o.$..... |
00000020 ef 0b f2 b7 c9 0f b5 e9 80 de 51 9a 16 f8 82 40 | .....Q....@|
00000030 f0 60 a8 d5 35 74 84 1c e5 97 43 XX XX XX XX XX | .`..5t....C....|
00000040 ff 5b 0c e9 07 8b d0 8e d2 fe a4 0e fb a5 ca 5d | .[.....] |
```

```
00000000 a4 07 9e f2 f5 57 6c 70 88 83 25 f0 86 11 8a e5 | .....Wlp..%..... |
00000010 6e 74 4e 61 6d 65 3e XX XX XX XX XX XX XX XX | ntName>XXXXXXXXXX |
00000020 34 3c 2f 41 63 63 6f 75 6e 74 4e 61 6d 65 3e 3c | 4</AccountName>< |
00000030 47 72 6f 75 70 73 3e 3c 49 44 3e XX XX XX XX XX | Groups><ID>XXXXXX |
00000040 3c 2f 49 44 3e 3c 4e 61 6d 65 3e 46 72 69 65 6e | </ID><Name>Frien |
```

We Have Both Kinds: AES and XOR (10)

```
00000000 11 11 11 11 11 11 11 11 11 11 11 11 11 11 11 | ..... |
00000010 11 11 11 11 11 11 11 11 11 11 11 11 11 11 11 | ..... |
00000020 ef 0b f2 b7 c9 0f b5 e9 80 de 51 9a 16 f8 82 40 | .....Q....@|
00000030 f0 60 a8 d5 35 74 84 1c e5 97 43 XX XX XX XX XX | .`..5t....C....|
00000040 ff 5b 0c e9 07 8b d0 8e d2 fe a4 0e fb a5 ca 5d | .[.....]|
```

```
00000000 a4 07 9e f2 f5 57 6c 70 88 83 25 f0 86 11 8a e5 | .....Wlp..%.....|
00000010 a4 07 9e f2 f5 57 6c 70 88 83 25 f0 86 11 8a e5 | .....Wlp..%.....|
00000020 34 3c 2f 41 63 63 6f 75 6e 74 4e 61 6d 65 3e 3c | 4</AccountName><|
00000030 47 72 6f 75 70 73 3e 3c 49 44 3e XX XX XX XX XX | Groups><ID>XXXXXX|
00000040 3c 2f 49 44 3e 3c 4e 61 6d 65 3e 46 72 69 65 6e | </ID><Name>Frien|
```

We Have Both Kinds: AES and XOR (11)

```
00000000 11 11 11 11 11 11 11 11 11 11 11 11 11 11 11 | ..... |
00000010 11 11 11 11 11 11 11 11 11 11 11 11 11 11 11 | ..... |
00000020 11 11 11 11 11 11 11 11 11 11 11 11 11 11 11 | ..... |
00000030 11 11 11 11 11 11 11 11 11 11 11 11 11 11 11 | ..... |
00000040 11 11 11 11 11 11 11 11 11 11 11 11 11 11 11 | ..... |
```

```
00000000 a4 07 9e f2 f5 57 6c 70 88 83 25 f0 86 11 8a e5 | .....Wlp..%..... |
00000010 a4 07 9e f2 f5 57 6c 70 88 83 25 f0 86 11 8a e5 | .....Wlp..%..... |
00000020 a4 07 9e f2 f5 57 6c 70 88 83 25 f0 86 11 8a e5 | .....Wlp..%..... |
00000030 a4 07 9e f2 f5 57 6c 70 88 83 25 f0 86 11 8a e5 | .....Wlp..%..... |
00000040 a4 07 9e f2 f5 57 6c 70 88 83 25 f0 86 11 8a e5 | .....Wlp..%..... |
```

We Have Both Kinds: AES and XOR (12)

- Did we mention that the DLL in question was deployed on both client and server components?
- This implied that any vulnerabilities in the DLL could potentially be leveraged on the server
 - Buffer overflows in the client-server traffic
 - Also, this gives an attacker an advantage since s/he has plenty of time to analyze the DLL offline

(Skip back)

We Have Both Kinds: AES and XOR (13)

- And now, the other problem: XOR
- The application allowed user credentials to be “remembered”
 - Stored in the Registry
 - Obscured with a position based XOR
 - Trivial to reverse
- What were they thinking?
- We've been down that road before... <sigh>

The House Always... Loses?

- Online gambling application

[\(Back\)](#) [\(Skip\)](#)

The House Always... Loses? (1)

- Online gambling application
 - Fat MS Windows client
 - Used a combination of web-based and home-grown protocols
 - Login credentials were hashed/obfuscated and stored in a local INF file
- Unfortunately, the credentials were not completely obfuscated and the scheme was cracked

The House Always... Loses? (2)

- It began with the username, which looked like a SHA1 (or any 160-bit) hash that had been specially formatted:
 - 35675E68-F4B5AF7B-995D9205-AD0FC438-42F16450
 - `^([0-9A-F]{8}-){4}[0-9A-F]{8}`
- A simple test (guest/guest) confirmed our theory:
 - `echo -n "guest" | openssl sha1 | perl -p -e 's/(.{8})/uc($1)."-"/ge; s/-$///;'`
- Using that as our basis, we conjectured that password obfuscation would follow a similar pattern

The House Always... Loses? (3)

- The obfuscated password hash looked like a SHA1 too, but unlike the username hash, it was Base64 encoded:
 - `VcqTUWCKPxpAhH8HBIL31kLxZFA=`
 - `^[+/0-9A-Za-z]{27}=`
- This time, a simple test did not confirm our theory:
 - `echo -n "guest" | openssl sha1 -binary | openssl base64`
 - `NwdeaPS1r3uZXZIFrQ/EOELxZFA=`
- However, it revealed something very interesting...

The House Always... Loses? (4)

- The low 32 bits of the obfuscated password hash were identical to the low 32 bits of the password's SHA1 hash
 - bdca9351608a3f1a40847f070482f7d642f16450
 - 35675e68f4b5af7b995d9205ad0fc43842f16450
- The obfuscation algorithms were recovered after several hours of disassembly
- From our perspective, it looked like they had upgraded from 128-bit to 160-bit hashes at some point (and shortcuts were taken)

The House Always... Loses? (5)

- Obfuscation routines were ported to C, which allowed us to produce our own credentials externally and much faster than going through the application
- Password cracking could have been done even without the disassembly effort
- Based on the amount of assembly analyzed, we conjecture that a significant amount of work went into obfuscating credentials – perhaps more than it took to de-obfuscate them

Can't Crack SSL? Just Talk Plaintext!

- A mobile application

[\(Back\)](#) [\(Skip\)](#)

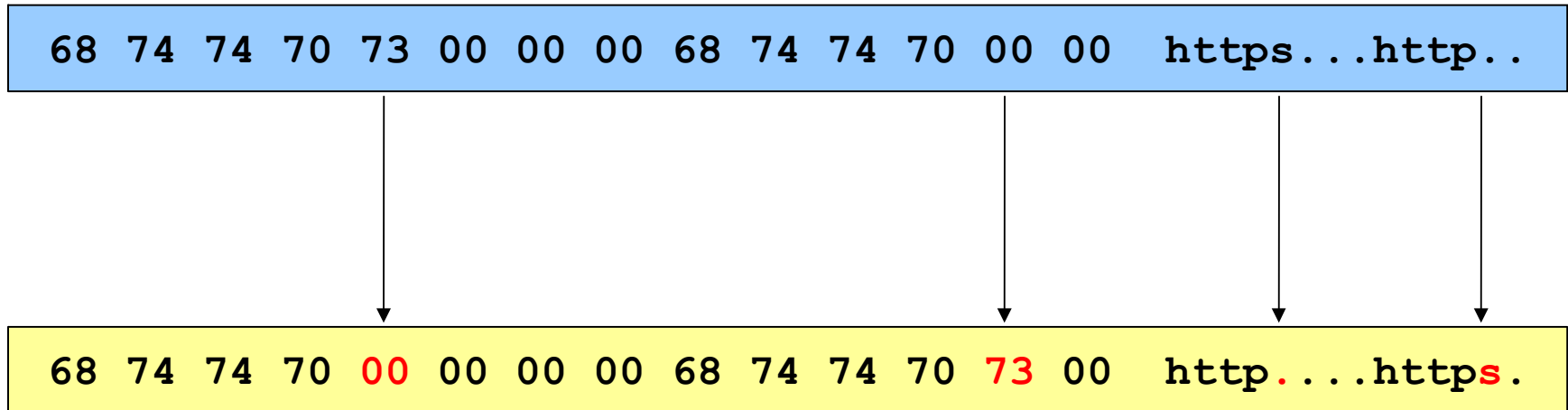
Can't Crack SSL? Just Talk Plaintext! (1)

- A mobile application
 - Fat BREW client that communicates via HTTPS
 - Strong/Encrypted authentication discouraged MITM attacks and eliminated visibility between handsets and servers
 - The provided handset was in test mode, so the application could be tweaked
 - NOTE: BREW applications are digitally signed, so this line of attack is normally cut off
- Suppose BREW protections could be bypassed...

Can't Crack SSL? Just Talk Plaintext! (2)

- We decided to find out how the system would react if the binary was modified to speak HTTP rather than HTTPS
 - Trivial binary edit – no assembly required
- The results were interesting
 - The handset happily spoke HTTP
 - The servers accepted the unencrypted traffic, and, most importantly, certificate authentication was gone!

Can't Crack SSL? Just Talk Plaintext! (3)



- The modification was a simple flip-flop of strings
- With the cloak removed, we were free to study and attack client-server communications at will
- MITM attacks were back on the menu simply by modifying URL addresses

Can't Crack SSL? Just Talk Plaintext! (4)

- The primary flaw, in our opinion, was the belief that the devices couldn't be hacked
- Designers should assume that all components of a system will be known by or become visible to the attacker sooner or later
 - We found leaked copies of the device flash images, and of tools that would put the devices into test mode, on public websites
- Designers should use every layer of protection to their advantage, but they should not assume that any one layer will suffice
 - Try to design each layer as if it were the only layer

Can't Crack SSL? Just Talk Plaintext! (5)

- If a production system doesn't require unencrypted communications, don't deploy production binaries and services that support unencrypted communications
- If a mixed HTTP/HTTPS webserver starts getting queries in plaintext that ought to be SSL-only... maybe that's a clue!

“Take my data. Please!”

- A compartmentalized manufacturing application

[\(Back\)](#) [\(Skip\)](#)

“Take my data. Please!” (1)

- A compartmentalized manufacturing application
 - Fat MS Windows client
 - Proxy account built into the client issues an SQL query to authenticate the user's login credentials and obtain the user's encrypted database credentials
 - Database traffic was not encrypted, but some of the individual fields were
- Unfortunately, this application used hard-coded keys, weak obfuscation techniques, and it leaked other users' encrypted credentials!

“Take my data. Please!” (2)

- The authentication process started like this:
 - User launches application (FOO.EXE)
 - User populates login form (user123/hammer)
 - Application issues the query shown below using hard-coded database credentials
 - u1, u2 = encrypted login username/password
 - u3, u4 = encrypted database username/password

```
SELECT id, u1, u2, u3, u4 FROM common_table WHERE somekey = 22;
```


“Take my data. Please!” (3)

- Next, the authentication process did this:
 - Decrypt u1 of each record until a match is found for the current user (user123)
 - Decrypt u2 to determine if password matches the one supplied by the current user (hammer)
 - If both u1 and u2 checks pass, decrypt u3 and u4 to obtain database credentials
 - At this point, the user could work on the project
- *All of these checks are done locally, in the client!*

“Take my data. Please!” (4)

- We could sniff the results of the query and could tell that the fields were encrypted
- We used OllyDbg to help us analyze the application
 - The encryption/decryption key was stored in a local file: “query.key”
 - 1536 bytes of binary data
 - After some custom code, DLL abuse, and disassembly, we found the key, but it was only 57 bytes long
 - So why was “query.key” so much bigger?

“Take my data. Please!” (5)

The screenshot shows OllyDbg debugging a process named 'd.exe'. The CPU window displays assembly instructions from address 10001D24 to 10001D70. The registers window shows the state of registers like EAX, ECX, and EDI. The hex dump window shows memory contents from address 001269E8 to 00126B68. A red circle highlights the hex value 'FE 16' at address 00126B58, with a red arrow pointing to the word 'KEY'.

Address	Hex dump	ASCII
001269E8	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
001269F8	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00126A08	40 06 13 00 00 00 00 00 00 00 00 00 00 00 00 00	@!.....
00126A18	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00126A28	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00126A38	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00126A48	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00126A58	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00126A68	90 2C 13 00 78 01 13 00 00 00 00 00 00 00 00 00	E!..00!.....
00126A78	90 2C 13 00 90 2C 13 00 88 2C 13 00 00 00 00 00	E!..00!..E!..
00126A88	00 00 00 00 00 00 00 00 90 06 00 00 00 00 00 00E!..E!..
00126A98	00 00 00 00 00 00 00 00 78 01 13 00 90 2C 13 0000!..E!..
00126AA8	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00126AB8	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00126AC8	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00126AD8	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00126AE8	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00126AF8	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00126B08	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00126B18	00 00 00 00 00 00 13 00 80 00 00 00 68 53 13 00!..hV!..
00126B28	68 53 13 00 00 00 00 00 68 01 13 00 47 00 00 00!..E!..G!..
00126B38	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00126B48	FE 16	..
00126B58		
00126B68		
00126B78	53 FB 77 08 06 13 00 05 00 00 04 6C 12 00 00	BU!..!..S!..*!..
00126B88	00 00 13 00 00 00 40 00 00 00 00 00 00 00 00 00
00126B98	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00126BA8	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00126BB8	00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00

“Take my data. Please!” (6)

- We recovered the algorithm and wrote an extraction program to pull the key...

```
...
#define INITIAL_OFFSET 0x232
#define STEP 7
#define SIZE 1536
#define KEY_LENGTH 0x39
int main() {
    ...
    iNRead = fread(pcData, 1, SIZE, pFile);
    ...
    for (i = 0, iOffset = INITIAL_OFFSET; i < KEY_LENGTH; i++, iOffset += STEP) {
        fprintf(stdout, "%02x,", pcData[iOffset]);
    }
    fprintf(stdout, "\n");
    return 0;
}
```

“Take my data. Please!” (7)

- Because multiple records were returned in the initial SQL query, we were able to decrypt other user's credentials
 - In fact, we were able to modify the query to select ALL records!
- We wrote a C program to abuse the exported DLL routines and the recovered key
 - LoadLibrary(), GetProcAddress(), and a bit of inline assembly
- Sure, we'll take your data... :)

- **Case Studies - Home Grown Crypto Wall of Shame**
 - Obfuscation Gone Bad
 - Keys? We Don't Need No Stinking Keys
 - We Don't Need Logic, We've Got Crypto!
 - Opening Your Neighbour's Garage Door
 - We Have Both Kinds: AES and XOR
 - The House Always... Loses?
 - Can't Crack SSL? Just Talk Plaintext!
 - “Take my data. Please!”

- Intro
- Case Studies - Home Grown Crypto Wall of Shame
- Prehash Analysis
- Recommendations and Closing Remarks

(Skip)

Prehash Analysis

- Session IDs are often created by combining several of the following variables:
 - Username, realm, IP address, device ID, resource, time, shared secrets, nonces, etc.
- Typically, several variables are concatenated and hashed with an algorithm such as MD5 or SHA1
- Too few, easily guessed, or badly ordered variables are common problems
 - Application context often makes all the difference

Prehash Analysis (2)

- Sometimes, an appropriate number of variables are chosen, but they're ordered in a way that makes the scheme weak
- Our MD5 prehash analysis confirmed that weak schemes could be attacked faster
- The idea was to find a way to decrease the amount of time required to brute force a particular scheme
- In particular, we focused on a class of IDs that combine two types of variables: fixed prefix and variable suffix

Prehash Analysis (3)

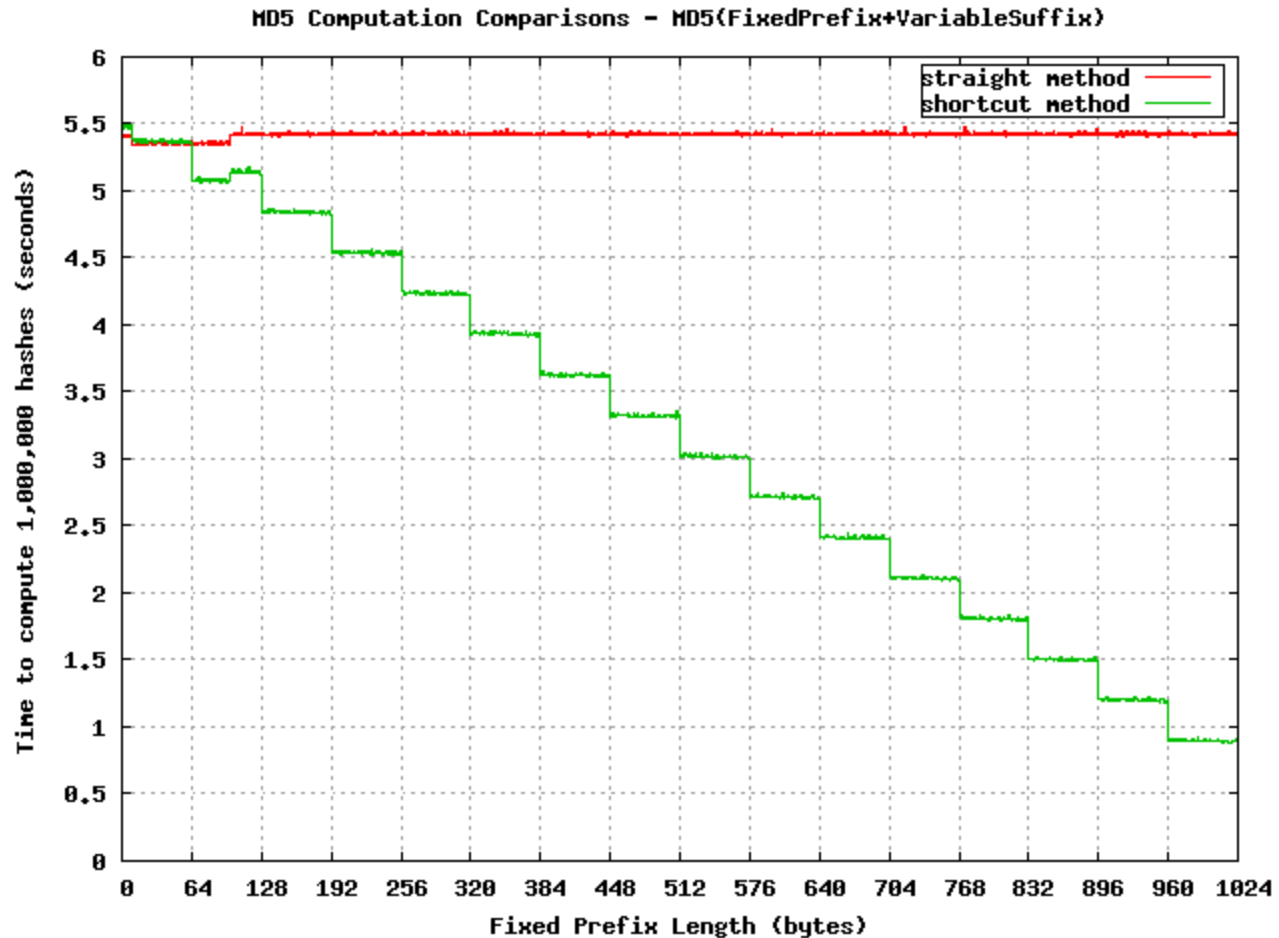


- Fixed Prefix = Easy to guess or infer
- Variable Suffix = Hard to guess or random
- Prehashing the fixed portion allows overall hashes to be computed faster
- If the variable portion is too short, this could be a real problem

Prehash Analysis (4)

- There is a significant advantage to prehashing if the hard-to-guess variable(s) are at the end of the string and the total string length is longer than 64 bytes
- At 512 bytes, the performance gain is roughly 2:1; at 1024, bytes the gain is roughly 6:1 (see graph)
- The results of our analysis are documented here:
 - www.korelogic.com/Resources/Papers/md5-analysis-prehash.txt
- Rainbow-table techniques should be applicable in some cases, for the same reason, making fixed-prefixes even worse

Prehash Analysis (5)



Prehash Analysis (6)

- Session ID: 1020-byte fixed prefix, 4-byte variable suffix
 - Prehash method will crack the ID in ~1 hour
 - Regular method will crack the ID in ~6.5 hours
 - Time required to try each ID is not included here
- Put the hard-to-guess stuff at or near the front
 - hash(unknown,known,known) – DO
 - hash(known,known,unknown) – DON'T
- Assume the easy-to-guess stuff will be lost
 - context, loose lips, lucky/educated guesses, etc.

Prehash Analysis (7)

- Don't make the hard-to-guess stuff too short
 - We've seen IP addresses, SSNs, timestamps, PIDs, phone numbers, etc. used as the “hard-to-guess” stuff
 - Too often, the space for these items can be reduced:
 - Targeting phone numbers in particular area code and exchange reduces the space from 10^{10} to 10^4
 - Targeting SSNs in Maryland (212-220) reduces the space from 10^9 to less than 10^7 (9×10^6); targeting specific years of birth can reduce it still further
 - Targeting IPs can be dead simple if the users all have a common ISP!
- Include brute force detection and logging in your designs

- Intro
- Case Studies - Home Grown Crypto Wall of Shame
- Prehash Analysis
- Recommendations and Closing Remarks

Don't Grow Your Own Crypto

- Unless you are a crypto mathematician and your job is to build crypto algorithms
- There are many ways to get it wrong...
- New algorithms require years of analysis by many experts to be verified and accepted
- Proprietary algorithms won't get the level scrutiny that open algorithms will
- Just because you can't break it, doesn't mean that someone else can't break it

Don't Assume Too Much

- “That'll never happen...” (epidemic)
- Plan on multiple protection mechanisms failing or being bypassed rather than assuming they'll just work
- “Don't underestimate the power of the dark side...”
 - Software disassembly techniques and tools are getting better all the time
 - If your secrets are important enough, someone will attempt to get them

Don't Rely On Obfuscation

- It's very difficult to prevent algorithm recovery when the attacker has full access to the client software, device, or whatever...
- Your time is better spent hardening server-side components and placing zero trust in client-side components that are out of your control
- Assume that any obfuscation you use is transparent
- Don't hide keys where they can be recovered – even if you think they can't be recovered

Design From Your Data Out

- Design security around your data and work outwards from there – not the other way 'round
- Use layers of protection whenever possible
 - Assume that each layer will be broken
 - Add detection and logging at every layer
- Keep protection mechanisms simple and modular
- Compartmentalize your data
 - Design such that a breach of one compartment does not automatically imply a breach of all

Design For Agility

- Crypto Agility
 - Hashes, keys, key lengths, algorithms, and RNGs
- Protocol Agility
 - How will your application handle protocol flaws?
- Error Agility
 - Does your application handle unexpected errors gracefully?
- Application Agility
 - How do you plan to upgrade all those clients?

Use Care When Implementing Crypto

- Know how and why you are using it
- Use standard libraries that have already worked out most of the kinks
- Use previously stated guiding principles
 - Don't grow your own crypto
 - Don't assume too much
 - Don't rely on obfuscation
 - Design from your data out
 - Design for agility

Hardware Trumps Software

- Generally, it costs more to reverse crypto and/or protection mechanisms implemented in hardware (if done well)
 - NOTE: We're not talking about appliances built around general purpose computers
- But don't get complacent... Hardware reversing is just a matter of time and money
 - What is the life expectancy of your product (e.g., client software, mobile devices, appliances)?
 - How much time and money do you want the attackers to spend attacking your product?

That's All Folks... Thanks!

- www.korelogic.com
- Random plug: <http://marc.info>
- Extra credit: Why did the fly dance on the jar?

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